

# $\omega_1$ under $\Pi_1$ -Collection

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## Abstract

We describe a proof-theoretic bound on  $\Sigma_2$ -definable countable ordinals in Kripke-Platek set theory with  $\Pi_1$ -Collection and the existence of  $\omega_1$ .

## 1 Introduction

Let  $(\omega_1)$  denote an axiom stating that ‘there exists an uncountable regular ordinal’, and  $T_1 := \text{KP}\omega + (V = L) + (\Pi_1\text{-Collection}) + (\omega_1)$ . Let  $\rho_0$  denote the least ordinal above  $\omega_1$  such that  $L_{\rho_0} \models (\Pi_1\text{-Collection})$ . In this note a collapsing function  $\Psi_{\omega_1} : \alpha \mapsto \Psi_{\omega_1}(\alpha) < \omega_1$  is defined, and it is shown that for each  $n < \omega$ ,  $T_1 \vdash \forall \alpha < \omega_n(\rho_0 + 1) \exists x < \omega_1 (x = \Psi_{\omega_1}(\alpha))$  with a  $\Sigma_2$ -formula  $x = \Psi_{\omega_1}(\alpha)$ , cf. Lemma 4.5. Conversely we show the

**Theorem 1.1** *For a sentence  $\exists x \in L_{\omega_1} \varphi(x)$  with a  $\Sigma_2$ -formula  $\varphi(x)$ , if*

$$T_1 \vdash \exists x \in L_{\omega_1} \varphi(x)$$

*then*

$$\exists n < \omega [T_1 \vdash \exists x \in L_{\Psi_{\omega_1}(\omega_n(\rho_0+1))} \varphi(x)].$$

This paper relies on our [1].

## 2 $\Sigma_1$ -Skolem hulls

Everything in this section is reproduced from [1].

For a model  $\langle M; \in \rangle (M \times M)$  and  $X \subset M$ ,  $\Sigma_1^M(X)$  denotes the set of  $\Sigma_1(X)$ -definable subsets of  $M$ , where  $\Sigma_1(X)$ -formulae may have parameters from  $X$ .  $\Sigma_1^M(M)$  is denoted  $\Sigma_1(M)$ .

An ordinal  $\alpha > 1$  is said to be a *multiplicative principal number* iff  $\alpha$  is closed under ordinal multiplication, i.e.,  $\exists \beta [\alpha = \omega^{\omega^\beta}]$ . If  $\alpha$  is a multiplicative principal number, then  $\alpha$  is closed under Gödel’s pairing function  $j$  and there exists a  $\Delta_1$ -bijection between  $\alpha$  and  $L_\alpha$  for the constructible hierarchy  $L_\alpha$  up to  $\alpha$ . In this section  $\sigma$  is assumed to be a multiplicative principal number  $> \omega$ .

**Definition 2.1** 1.  $cf(\kappa) := \min\{\alpha \leq \kappa : \text{there is a cofinal map } f : \alpha \rightarrow \kappa\}$ .

2.  $\rho(L_\sigma)$  denotes the  $\Sigma_1$ -projectum of  $L_\sigma$ :  $\rho(L_\sigma)$  is the least ordinal  $\rho$  such that  $\mathcal{P}(\rho) \cap \Sigma_1(L_\sigma) \not\subset L_\sigma$ .

3. Let  $\alpha \leq \beta$  and  $f : L_\alpha \rightarrow L_\beta$ . Then the map  $f$  is a  $\Sigma_1$ -elementary embedding, denoted  $f : L_\alpha \prec_{\Sigma_1} L_\beta$  iff for any  $\Sigma_1(L_\alpha)$ -sentence  $\varphi[\bar{a}]$  ( $\bar{a} \subset L_\alpha$ ),  $L_\alpha \models \varphi[\bar{a}] \Leftrightarrow L_\beta \models \varphi[f(\bar{a})]$  where  $f(\bar{a}) = f(a_1), \dots, f(a_k)$  for  $\bar{a} = a_1, \dots, a_k$ . An ordinal  $\gamma$  such that  $\forall \delta < \gamma [f(\delta) = \delta] \& f(\gamma) > \gamma$  is said to be the *critical point* of the  $\Sigma_1$ -elementary embedding  $f$  if such an ordinal  $\gamma$  exists.

4. For  $X \subset L_\sigma$ ,  $\text{Hull}_{\Sigma_1}^\sigma(X)$  denotes the set ( $\Sigma_1$ -Skolem hull of  $X$  in  $L_\sigma$ ) defined as follows.  $<_L$  denotes a  $\Delta_1$ -well ordering of the constructible universe  $L$ . Let  $\{\varphi_i : i \in \omega\}$  denote an enumeration of  $\Sigma_1$ -formulae in the language  $\{\in\}$ . Each is of the form  $\varphi_i \equiv \exists y \theta_i(x, y; u)$  ( $\theta \in \Delta_0$ ) with fixed variables  $x, y, u$ . Set for  $b \in X$

$$\begin{aligned} r_{\Sigma_1}^\sigma(i, b) &\simeq \text{the } <_L \text{-least } c \in L_\sigma \text{ such that } L_\sigma \models \theta_i((c)_0, (c)_1; b) \\ h_{\Sigma_1}^\sigma(i, b) &\simeq (r_{\Sigma_1}^\sigma(i, b))_0 \\ \text{Hull}_{\Sigma_1}^\sigma(X) &= \text{rng}(h_{\Sigma_1}^\sigma) = \{h_{\Sigma_1}^\sigma(i, b) \in L_\sigma : i \in \omega, b \in X\} \end{aligned}$$

Then  $L_\sigma \models \exists x \exists y \theta_i(x, y; b) \rightarrow h_{\Sigma_1}^\sigma(i, b) \downarrow \& \exists y \theta_i(h_{\Sigma_1}^\sigma(i, b), y; b)$ .

**Proposition 2.2** Assume that  $X$  is a set in  $L_\sigma$ . Then  $r_{\Sigma_1}^\sigma$  and  $h_{\Sigma_1}^\sigma$  are partial  $\Delta_1(L_\sigma)$ -maps such that the domain of  $h_{\Sigma_1}^\sigma$  is a  $\Sigma_1(L_\sigma)$ -subset of  $\omega \times X$ . Therefore its range  $\text{Hull}_{\Sigma_1}^\sigma(X)$  is a  $\Sigma_1(L_\sigma)$ -subset of  $L_\sigma$ .

**Proposition 2.3** Let  $Y = \text{Hull}_{\Sigma_1}^\sigma(X)$ . For any  $\Sigma_1(Y)$ -sentence  $\varphi(\bar{a})$  with parameters  $\bar{a}$  from  $Y$   $L_\sigma \models \varphi(\bar{a}) \Leftrightarrow Y \models \varphi(\bar{a})$ . Namely  $Y \prec_{\Sigma_1} L_\sigma$ .

**Definition 2.4** (Mostowski collapsing function  $F$ )

By Proposition 2.3 and the Condensation Lemma we have an isomorphism (Mostowski collapsing function)

$$F : \text{Hull}_{\Sigma_1}^\sigma(X) \leftrightarrow L_\gamma$$

for an ordinal  $\gamma \leq \sigma$  such that  $F \upharpoonright Y = id \upharpoonright Y$  for any transitive  $Y \subset \text{Hull}_{\Sigma_1}^\sigma(X)$ .

Let us denote, though  $\sigma \notin \text{dom}(F) = \text{Hull}_{\Sigma_1}^\sigma(X)$

$$F(\sigma) := \gamma.$$

Also for the above Mostowski collapsing map  $F$  let

$$F^{\Sigma_1}(x; \sigma, X) := F(x).$$

The inverse  $G := F^{-1}$  of  $F$  is a  $\Sigma_1$ -elementary embedding from  $L_{F(\sigma)}$  to  $L_\sigma$ .

**Proposition 2.5** *Let  $L_\sigma \models \text{KP}\omega + \Sigma_1\text{-Collection}$ . Then for  $\kappa \leq \sigma$ ,  $\{(x, y) : x < \kappa \ \& \ y = \min\{y < \kappa : \text{Hull}_{\Sigma_1}^\sigma(x \cup \{\kappa\}) \cap \kappa \subset y\}\}$  is a  $\text{Bool}(\Sigma_1(L_\sigma))$ -predicate on  $\kappa$ , and hence a set in  $L_\sigma$  if  $\kappa < \sigma$  and  $L_\sigma \models \Sigma_1\text{-Separation}$ .*

$F_{x \cup \{\kappa\}}^{\Sigma_1}(y)$  denotes the Mostowski collapse  $F^{\Sigma_1}(y; \sigma, x \cup \{\kappa\})$ .

**Theorem 2.6** *Let  $\sigma$  be an ordinal such that  $L_\sigma \models \text{KP}\omega + \Sigma_1\text{-Separation}$ , and  $\omega \leq \alpha < \kappa < \sigma$  with  $\alpha$  a multiplicative principal number and  $\kappa$  a limit ordinal. Then the following conditions are mutually equivalent:*

1.  $L_\sigma \models {}^\alpha \kappa \subset L_\kappa$ .
2.  $L_\sigma \models \alpha < cf(\kappa)$ .
3. *There exists an ordinal  $x$  such that  $\alpha < x < \kappa$ ,  $\text{Hull}_{\Sigma_1}^\sigma(x \cup \{\kappa\}) \cap \kappa \subset x$  and  $F_{x \cup \{\kappa\}}^{\Sigma_1}(\sigma) < \kappa$ .*
4. *For the Mostowski collapse  $F_{x \cup \{\kappa\}}^{\Sigma_1}(y)$ , there exists an ordinal  $x$  such that  $\alpha < x = F_{x \cup \{\kappa\}}^{\Sigma_1}(\kappa) < F_{x \cup \{\kappa\}}^{\Sigma_1}(\sigma) < \kappa$ , and for any  $\Sigma_1$ -formula  $\varphi$  and any  $a \in L_x$ ,  $L_\sigma \models \varphi[\kappa, a] \rightarrow L_{F_{x \cup \{\kappa\}}^{\Sigma_1}(\sigma)} \models \varphi[x, a]$  holds.*

**Definition 2.7**  $T_1 := \text{KP}\omega + (V = L) + (\Pi_1\text{-Collection}) + (\omega_1)$  denotes an extension of the Kripke-Platek set theory with the axioms of infinity, constructibility,  $\Pi_1$ -Collection and the following axiom:

$$(\omega_1) \exists \kappa \forall \alpha < \kappa \exists \beta, \gamma < \kappa [\alpha < \beta < \gamma \wedge L_\gamma = \text{rng}(F_{\beta \cup \{\kappa\}}^{\Sigma_1}) \wedge \text{Hull}_{\Sigma_1}(\beta \cup \{\kappa\}) \cap \kappa \subset \beta]$$

where  $F_{\beta \cup \{\kappa\}}^{\Sigma_1} : \text{Hull}_{\Sigma_1}(\beta \cup \{\kappa\}) \rightarrow L_\gamma$  is the Mostowski collapsing map, and  $\text{Hull}_{\Sigma_1}(x)$  is the  $\Sigma_1$ -Skolem hull of sets  $x$  in the universe.

From Theorem 2.6 we see that  $T_1 \vdash \exists \kappa \forall \alpha < \kappa (\alpha < cf(\kappa))$ .

### 3 A theory equivalent to $T_1$

Referring Theorem 2.6 let us interpret  $T_1$  to another theory. The base language here is  $\{\in\}$ .

Let  $\rho_0$  denotes the least ordinal above the least uncountable ordinal  $\omega_1$  such that  $L_{\rho_0} \models (\Pi_1\text{-Collection})$ .  $F_X(x) := F^{\Sigma_1}(x; \rho_0, X)$  and  $\text{Hull}(X) := \text{Hull}_{\Sigma_1}^{\rho_0}(X)$ .

The predicate  $P$  is intended to denote the relation  $P(x, y)$  iff  $x = F_{x \cup \{\omega_1\}}(\omega_1)$  and  $y = F_{x \cup \{\omega_1\}}(\rho_0)$ . Also the predicate  $P_{\rho_0}(x)$  is intended to denote the relation  $P_{\rho_0}(x)$  iff  $x = F_x(\rho_0)$ .

**Definition 3.1**  $T(\omega_1)$  denotes the set theory defined as follows.

1. Its language is  $\{\in, P, P_{\rho_0}, \omega_1\}$  for a binary predicate  $P$ , a unary predicate  $P_{\rho_0}$  and an individual constant  $\omega_1$ .

2. Its axioms are obtained from those of  $\mathbf{KP}\omega + (\Pi_1\text{-Collection})$  in the expanded language <sup>1</sup>, the axiom of constructibility  $V = L$  together with the axiom schema saying that  $\omega_1$  is an uncountable regular ordinal, cf. (2) and (1), and if  $P(x, y)$  then  $x$  is a critical point of the  $\Sigma_1$ -elementary embedding from  $L_y \cong \text{Hull}(x \cup \{\omega_1\})$  to the universe  $L_{\rho_0}$ , cf. (1), and if  $P_{\rho_0}(x)$  then  $x$  is a critical point of the  $\Sigma_1$ -elementary embedding from  $L_x \cong \text{Hull}(x)$  to the universe  $L_{\rho_0}$ , cf. (3): for a formula  $\varphi$  and an ordinal  $\alpha$ ,  $\varphi^\alpha$  denotes the result of restricting every unbounded quantifier  $\exists z, \forall z$  in  $\varphi$  to  $\exists z \in L_\alpha, \forall z \in L_\alpha$ .

- (a)  $x \in \text{Ord}$  is a  $\Delta_0$ -formula saying that ‘ $x$  is an ordinal’.  
 $(\omega < \omega_1 \in \text{Ord}), (P(x, y) \rightarrow \{x, y\} \subset \text{Ord} \wedge x < y < \omega_1)$  and  
 $(P_{\rho_0}(x) \rightarrow x \in \text{Ord})$ .

(b)

$$P(x, y) \rightarrow a \in L_x \rightarrow \varphi[\omega_1, a] \rightarrow \varphi^y[x, a] \quad (1)$$

for any  $\Sigma_1$ -formula  $\varphi$  in the language  $\{\in\}$ .

(c)

$$a \in \text{Ord} \cap \omega_1 \rightarrow \exists x, y \in \text{Ord} \cap \omega_1 [a < x \wedge P(x, y)] \quad (2)$$

(d)

$$P_{\rho_0}(x) \rightarrow a \in L_x \rightarrow \varphi[a] \rightarrow \varphi^x[a] \quad (3)$$

for any  $\Sigma_1$ -formula  $\varphi$  in the language  $\{\in\}$ .

(e)

$$a \in \text{Ord} \rightarrow \exists x \in \text{Ord} [a < x \wedge P_{\rho_0}(x)] \quad (4)$$

**Remark.** Though the axioms (3) and (4) for the  $\Pi_1$ -definable predicate  $P_{\rho_0}(x)$  are derivable from  $\Pi_1$ -Collection, the primitive predicate symbol  $P_{\rho_0}(x)$  is useful for our prof-theoretic study, cf. the proof of Lemma 5.20 below.

**Lemma 3.2**  $T(\omega_1)$  is a conservative extension of the set theory  $T_1$ .

**Proof.** First consider the axioms of  $T_1$  in  $T(\omega_1)$ . The axiom  $(\omega_1)$  follows from (1). Hence we have shown that  $T_1$  is contained in  $T(\omega_1)$ .

Next we show that  $T(\omega_1)$  is interpretable in  $T_1$ . Let  $\kappa$  be an ordinal in the axiom  $(\omega_1)$ . Interpret the predicate  $P(x, y) \leftrightarrow \{x, y\} \subset \text{Ord} \wedge (\text{Hull}(x \cup \{\kappa\}) \cap \kappa \subset x) \wedge (y = \sup\{F_{x \cup \{\kappa\}}(a) : a \in \text{Hull}(x \cup \{\kappa\})\})$ . We see from Theorem 2.6 that the interpreted (1) and (2) are provable in  $T_1$ .

It remains to show the interpreted (3) and (4) in  $T_1$ . It suffices to show that given an ordinal  $\alpha$ , there exists an ordinal  $x > \alpha$  such that  $\text{Hull}(x) \cap \text{Ord} \subset x$ .

First we show that for any  $\alpha$  there exists a  $\beta$  such that  $\text{Hull}(\alpha) \cap \text{Ord} \subset \beta$ . By Proposition 2.2 let  $h_{\Sigma_1}^{\rho_0}$  be the  $\Delta_1$ -surjection from the  $\Sigma_1$ -subset  $\text{dom}(h_{\Sigma_1}^{\rho_0})$  of

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<sup>1</sup> This means that the predicates  $P, P_{\rho_0}$  do not occur in  $\Delta_0$ -formulae for  $\Delta_0$ -Separation and  $\Pi_1$ -formulae  $\Pi_1$ -Collection.

$\omega \times \alpha$  to  $\text{Hull}(\alpha)$ , which is a  $\Sigma_1$ -class. From  $\Sigma_1$ -Separation we see that  $\text{dom}(h_{\Sigma_1}^{\rho_0})$  is a set. Hence by  $\Sigma_1$ -Collection,  $\text{Hull}(\alpha) = \text{rng}(h_{\Sigma_1}^{\rho_0})$  is a set. Therefore the ordinal  $\sup(\text{Hull}(\alpha) \cap \text{Ord})$  exists in the universe.

As in Proposition 2.5 we see that  $X = \{(\alpha, \beta) : \beta = \min\{\beta \in \text{Ord} : \text{Hull}(\alpha) \cap \text{Ord} \subset \beta\}\}$  is a set in  $L_{\rho_0}$  as follows. Let  $\varphi(\beta)$  be the  $\Pi_1$ -predicate  $\varphi(\beta) : \Leftrightarrow \forall \gamma \in \text{Ord}[\gamma \in \text{Hull}(\alpha) \rightarrow \gamma \in \beta]$ . Then  $\beta = \min\{\beta : \text{Hull}(\alpha) \cap \text{Ord} \subset \beta\}$  iff  $\varphi(\beta) \wedge \forall \gamma < \beta \neg \varphi(\gamma)$ , which is  $\text{Bool}(\Sigma_1(L_{\rho_0}))$  by  $\Pi_0$ -Collection. Hence  $X$  is a set in  $L_{\rho_0}$ .

Define recursively ordinals  $\{x_n\}_n$  as follows.  $x_0 = \alpha + 1$ , and  $x_{n+1}$  is defined to be the least ordinal  $x_{n+1}$  such that  $\text{Hull}(x_n) \cap \text{Ord} \subset x_{n+1}$ , i.e.,  $(x_n, x_{n+1}) \in X$ . We see inductively that such an ordinal  $x_n$  exists. Moreover  $n \mapsto x_n$  is a  $\Delta_1$ -map. Then  $x = \sup_n x_n < \rho_0$  is a desired one.  $\square$

## 4 Ordinals for $\omega_1$

For our proof-theoretic analysis of  $T_1$ , we need to talk about ‘ordinals’ less than the next epsilon number to the order type of the class of ordinals inside  $T_1$ . Let  $\text{Ord}^\varepsilon \subset V$  and  $<^\varepsilon$  be  $\Delta$ -predicates such that for any transitive and wellfounded model  $V$  of  $\text{KP}\omega$ ,  $<^\varepsilon$  is a well ordering of type  $\varepsilon_{\rho_0+1}$  on  $\text{Ord}^\varepsilon$  for the order type  $\rho_0$  of the class  $\text{Ord}$  in  $V$ .  $<^\varepsilon$  is seen to be a canonical ordering as stated in the following Proposition 4.1.

**Proposition 4.1** 1.  $\text{KP}\omega$  proves the fact that  $<^\varepsilon$  is a linear ordering.

2. For any formula  $\varphi$  and each  $n < \omega$ ,

$$\text{KP}\omega \vdash \forall x \in \text{Ord}^\varepsilon (\forall y <^\varepsilon x \varphi(y) \rightarrow \varphi(x)) \rightarrow \forall x <^\varepsilon \omega_n(\rho_0 + 1) \varphi(x) \quad (5)$$

In what follows of this section we work in  $T_1$ . For simplicity let us identify the code  $x \in \text{Ord}^\varepsilon$  with the ‘ordinal’ coded by  $x$ , and  $<^\varepsilon$  is denoted by  $<$  when no confusion likely occurs. Note that the ordinal  $\rho_0$  is the order type of the class of ordinals in the intended model  $L_{\rho_0}$  of  $T_1$ . Define simultaneously the classes  $\mathcal{H}_\alpha(X) \subset L_{\rho_0} \cup \varepsilon_{\rho_0+1}$  and the ordinals  $\Psi_{\omega_1}(\alpha)$  and  $\Psi_{\rho_0}(\alpha)$  for  $\alpha <^\varepsilon \varepsilon_{\rho_0+1}$  and sets  $X \subset L_{\omega_1}$  as follows. We see that  $\mathcal{H}_\alpha(X)$  and  $\Psi_\kappa(\alpha)$  ( $\kappa \in \{\omega_1, \rho_0\}$ ) are (first-order) definable as a fixed point in  $T_1$ , cf. Proposition 4.4.

Recall that  $\text{Hull}(X) = \text{Hull}_{\Sigma_1}^{\rho_0}(X) \subset L_{\rho_0}$  and  $F_X(x) = F^{\Sigma_1}(x; \rho_0, X)$  with  $F_X : \text{Hull}(X) \rightarrow L_\gamma$  for  $X \subset L_{\rho_0}$  and a  $F_X(\rho_0) = \gamma \leq \rho_0$ .

**Definition 4.2**  $\mathcal{H}_\alpha(X)$  is the Skolem hull of  $\{0, \omega_1, \rho_0\} \cup X$  under the functions  $+, \alpha \mapsto \omega^\alpha, \Psi_{\omega_1} \upharpoonright \alpha, \Psi_{\rho_0} \upharpoonright \alpha$ , the  $\Sigma_1$ -definability, and the Mostowski collapsing functions  $(x, d) \mapsto F_{x \cup \{\omega_1\}}(d)$  ( $\text{Hull}(x \cup \{\omega_1\}) \cap \omega_1 \subset x$ ) and  $d \mapsto F_x(d)$  ( $\text{Hull}(x) \cap \rho_0 \subset x$ ).

1.  $\{0, \omega_1, \rho_0\} \cup X \subset \mathcal{H}_\alpha(X)$ .
2.  $x, y \in \mathcal{H}_\alpha(X) \Rightarrow x + y, \omega^x \in \mathcal{H}_\alpha(X)$ .
3.  $\gamma \in \mathcal{H}_\alpha(X) \cap \alpha \Rightarrow \Psi_\kappa(\gamma) \in \mathcal{H}_\alpha(X)$  for  $\kappa \in \{\omega_1, \rho_0\}$ .

4.  $\text{Hull}(\mathcal{H}_\alpha(X) \cap L_{\rho_0}) \subset \mathcal{H}_\alpha(X)$ .

Namely for any  $\Sigma_1$ -formula  $\varphi[x, \vec{y}]$  in the language  $\{\in\}$  and parameters  $\vec{a} \subset \mathcal{H}_\alpha(X) \cap L_{\rho_0}$ , if  $b \in L_{\rho_0}$ ,  $(L_{\rho_0}, \in) \models \varphi[b, \vec{a}]$  and  $(L_{\rho_0}, \in) \models \exists! x \varphi[x, \vec{a}]$ , then  $b \in \mathcal{H}_\alpha(X)$ .

5. If  $x \in \mathcal{H}_\alpha(X) \cap \omega_1$  with  $\text{Hull}(x \cup \{\omega_1\}) \cap \omega_1 \subset x$ , and  $d \in (\text{Hull}(x \cup \{\omega_1\}) \cap \mathcal{H}_\alpha(X)) \cup \{\rho_0\}$ , then  $F_{x \cup \{\omega_1\}}(d) \in \mathcal{H}_\alpha(X)$ .

6. If  $x \in \mathcal{H}_\alpha(X) \cap \rho_0$  with  $\text{Hull}(x) \cap \rho_0 \subset x$ , and  $d \in (\text{Hull}(x) \cap \mathcal{H}_\alpha(X)) \cup \{\rho_0\}$ , then  $F_x(d) \in \mathcal{H}_\alpha(X)$ .

For  $\kappa \in \{\omega_1, \rho_0\}$

$$\Psi_\kappa(\alpha) := \min\{\beta \leq \kappa : \mathcal{H}_\alpha(\beta) \cap \kappa \subset \beta\}.$$

The ordinal  $\Psi_\kappa(\alpha)$  is well defined and  $\Psi_\kappa(\alpha) \leq \kappa$  for  $\kappa \in \{\omega_1, \rho_0\}$ .

**Proposition 4.3** 1.  $\mathcal{H}_\alpha(X)$  is closed under  $\Sigma_1$ -definability:  $\vec{a} \subset \mathcal{H}_\alpha(X) \cap L_{\rho_0} \Rightarrow \text{Hull}(\vec{a}) \subset \mathcal{H}_\alpha(X)$ .

2.  $\text{Hull}(\Psi_{\omega_1}(\alpha) \cup \{\omega_1\}) \cap \omega_1 = \Psi_{\omega_1}(\alpha)$  and  $\text{Hull}(\Psi_{\rho_0}(\alpha)) \cap \rho_0 = \Psi_{\rho_0}(\alpha) > \omega_1$

3.  $\mathcal{H}_\alpha(X)$  is closed under the Veblen function  $\varphi$  on  $\rho_0$ ,  $x, y \in \mathcal{H}_\alpha(X) \cap \rho_0 \Rightarrow \varphi xy \in \mathcal{H}_\alpha(X)$ .

4. If  $x \in \mathcal{H}_\alpha(X) \cap \omega_1$ ,  $\text{Hull}(x \cup \{\omega_1\}) \cap \omega_1 \subset x$ , and  $\delta \in (\text{Hull}(x \cup \{\omega_1\}) \cap \mathcal{H}_\alpha(X)) \cup \{\rho_0\}$ , then  $F_{x \cup \{\omega_1\}}(\delta) \in \mathcal{H}_\alpha(X)$ .

5. If  $x \in \mathcal{H}_\alpha(X) \cap \rho_0$ ,  $\text{Hull}(x) \cap \rho_0 \subset x$ , and  $\delta \in (\text{Hull}(x) \cap \mathcal{H}_\alpha(X)) \cup \{\rho_0\}$ , then  $F_x(\delta) \in \mathcal{H}_\alpha(X)$ .

The following Proposition 4.4 is easy to see.

**Proposition 4.4** Both of  $x = \mathcal{H}_\alpha(X)$  and  $y = \Psi_\kappa(\alpha)$  ( $\kappa \in \{\omega_1, \rho_0\}$ ) are  $\Sigma_2$ -predicates as fixed points in  $\text{KP}\omega$ .

**Lemma 4.5** For each  $n < \omega$ ,

$$T_1 \vdash \forall \alpha < \omega_{n+1}(\rho_0 + 1) \forall \kappa \in \{\omega_1, \rho_0\} \exists x < \kappa [x = \Psi_\kappa(\alpha)].$$

**Proof.** Let  $\kappa \in \{\omega_1, \rho_0\}$ . By Proposition 4.4 both  $x = \mathcal{H}_\alpha(\beta)$  and  $y = \Psi_\kappa(\alpha)$  are  $\Sigma_2$ -predicates. We show that  $A(\alpha) :\Leftrightarrow \forall \beta < \rho_0 \exists x [x = \mathcal{H}_\alpha(\beta)] \wedge \forall \kappa \in \{\omega_1, \rho_0\} \exists \beta < \kappa [\Psi_\kappa(\alpha) = \beta]$  is progressive. Then  $\forall \alpha < \omega_{n+1}(\rho_0 + 1) \forall \kappa \in \{\omega_1, \rho_0\} \exists x < \kappa [x = \Psi_\kappa(\alpha)]$  will follow from transfinite induction up to  $\omega_{n+1}(\rho_0 + 1)$ , cf. (5) in Proposition 4.1.

Assume  $\forall \gamma < \alpha A(\gamma)$  as our IH. Since  $\text{dom}(h_{\Sigma_1}^{\rho_0})$  is a  $\Sigma_1$ -subset of  $\omega \times \beta$  for  $\beta < \rho_0$ , it is a set by  $\Sigma_1$ -Separation. Then so is the image  $\text{Hull}(\beta)$  of the  $\Delta_1$ -map  $h_{\Sigma_1}^{\rho_0}$ . Hence  $\forall \beta < \rho_0 \exists h [h = \text{Hull}(\beta)]$ .

We see from this, IH and  $\Sigma_2$ -Collection that  $\forall \beta < \rho_0 \exists x [x = \mathcal{H}_\alpha(\beta) = \bigcup_m \mathcal{H}_\alpha^m(\beta)]$ , where  $\mathcal{H}_\alpha^m(\beta)$  is an  $m$ -th stage of the construction of  $\mathcal{H}_\alpha(\beta)$  such that  $x = \mathcal{H}_\alpha^m(\beta)$  is a  $\Sigma_2$ -predicate.

Define recursively ordinals  $\{\beta_m\}_m$  for  $\kappa \in \{\omega_1, \rho_0\}$  as follows.  $\beta_0 = 0$ , and  $\beta_{m+1}$  is defined to be the least ordinal  $\beta_{m+1} \leq \kappa$  such that  $\mathcal{H}_\alpha(\beta_m) \cap \kappa \subset \beta_{m+1}$ .

We see inductively that  $\beta_m < \kappa$  using the regularity of  $\kappa$  and the facts that  $\forall \beta < \kappa \exists x [x = \mathcal{H}_\alpha(\beta) \wedge \text{card}(x) < \kappa]$ , where  $\text{card}(x) < \kappa$  designates that there exists a surjection  $f : \gamma \rightarrow x$  for a  $\gamma < \kappa$  and  $f \in L_{\rho_0}$ . Moreover  $m \mapsto \beta_m$  is a  $\Sigma_2$ -map. Therefore  $\beta = \sup_m \beta_m < \kappa$  enjoys  $\mathcal{H}_\alpha(\beta) \cap \kappa \subset \beta$ .  $\square$

## 5 Operator controlled derivations for $T_1$

### 5.1 An intuitionistic fixed point theory $\text{FiX}^i(T_1)$

Let us introduce an intuitionistic fixed point theory  $\text{FiX}^i(T_1)$  over the set theory  $T_1$ . Fix an  $X$ -strictly positive formula  $\mathcal{Q}(X, x)$  in the language  $\{\in, =, X\}$  with an extra unary predicate symbol  $X$ . In  $\mathcal{Q}(X, x)$  the predicate symbol  $X$  occurs only strictly positive. The language of  $\text{FiX}^i(T_1)$  is  $\{\in, =, Q\}$  with a fresh unary predicate symbol  $Q$ . The axioms in  $\text{FiX}^i(T_1)$  consist of the following:

1. All derivable sentences in  $T_1$  in the language  $\{\in\}$ .
2. Induction schema for any formula  $\varphi$  in  $\{\in, =, Q\}$ :  
 $\forall x (\forall y \in x \varphi(y) \rightarrow \varphi(x)) \rightarrow \forall x \varphi(x)$ .
3. Fixed point axiom:  $\forall x [Q(x) \leftrightarrow \mathcal{Q}(Q, x)]$ .

The underlying logic in  $\text{FiX}^i(T_1)$  is defined to be the intuitionistic first-order predicate logic with equality.

**Lemma 5.1** *Let  $<^\varepsilon$  denote a  $\Delta_1$ -predicate defined in section 4. For each  $n < \omega$  and each formula  $\varphi$  in  $\{\in, =, Q\}$ ,*

$$\text{FiX}^i(T_1) \vdash \forall x (\forall y <^\varepsilon x \varphi(y) \rightarrow \varphi(x)) \rightarrow \forall x <^\varepsilon \omega_n(\rho_0 + 1) \varphi(x).$$

The following Theorem 5.2 is shown in [2].

**Theorem 5.2**  *$\text{FiX}^i(T_1)$  is a conservative extension of  $T_1$ .*

### 5.2 Classes of formulae

In this subsection we work in  $T_1$ .

The language  $\mathcal{L}_c$  is obtained from  $\{\in, P, P_{\rho_0}, \omega_1\}$  by adding names (individual constants)  $c_a$  of each set  $a \in L_{\rho_0}$ .  $c_a$  is identified with  $a$ . A *term* in  $\mathcal{L}_c$  is either a variable or a constant in  $L_{\rho_0}$ . Formulae in this language are defined in the next definition. Formulae are assumed to be in negation normal form.

**Definition 5.3** 1. Let  $t_1, \dots, t_m$  be terms. For each  $m$ -ary predicate constant  $R \in \{\in, P, P_{\rho_0}\}$   $R(t_1, \dots, t_m)$  and  $\neg R(t_1, \dots, t_m)$  are formulae, where  $m = 1, 2$ . These are called *literals*.

2. If  $A$  and  $B$  are formulae, then so are  $A \wedge B$  and  $A \vee B$ .
3. Let  $t$  be a term. If  $A$  is a formula and the variable  $x$  does not occur in  $t$ , then  $\exists x \in t A$  and  $\forall x \in t A$  are formulae.  $\exists x \in t A$ ,  $\forall x \in t A$  are *bounded quantifiers*.
4. If  $A$  is a formula and  $x$  a variable, then  $\exists x A$  and  $\forall x A$  are formulae. Unbounded quantifiers  $\exists x, \forall x$  are denoted by  $\exists x \in L_{\rho_0}, \forall x \in L_{\rho_0}$ , resp.

For formulae  $A$  in  $\mathcal{L}_c$ ,  $\mathbf{qk}(A)$  denotes the finite set of sets  $a \in L_{\rho_0}$  which are bounds of bounded quantifiers  $\exists x \in a, \forall x \in a$  occurring in  $A$ . Moreover  $\mathbf{k}(A)$  denotes the set of sets occurring in  $A$ .  $\mathbf{k}(A)$  is defined to include bounds of bounded quantifiers. By definition we set  $0 \in \mathbf{qk}(A)$ . Thus  $0 \in \mathbf{qk}(A) \subset \mathbf{k}(A) \subset L_{\rho_0}$ .

**Definition 5.4** 1.  $\mathbf{k}(\neg A) = \mathbf{k}(A)$  and similarly for  $\mathbf{qk}$ .

2.  $\mathbf{qk}(M) = \{0\}$  for any literal  $M$ .
3.  $\mathbf{k}(Q(t_1, \dots, t_m)) = (\{t_1, \dots, t_m\} \cap L_{\rho_0}) \cup \{0\}$  for literals  $Q(t_1, \dots, t_m)$  with predicates  $Q$  in the set  $\{\in, P, P_{\rho_0}\}$ .
4.  $\mathbf{k}(A_0 \vee A_1) = \mathbf{k}(A_0) \cup \mathbf{k}(A_1)$  and similarly for  $\mathbf{qk}$ .
5. For unbounded quantifiers,  $\mathbf{k}(\exists x A(x)) = \mathbf{k}(A(x))$  and similarly for  $\mathbf{qk}$ .
6. For bounded quantifiers with  $a \in L_{\rho_0}$ ,  $\mathbf{k}(\exists x \in a A(x)) = \{a\} \cup \mathbf{k}(A(x))$  and similarly for  $\mathbf{qk}$ .
7. For variables  $y$ ,  $\mathbf{k}(\exists x \in y A(x)) = \mathbf{k}(A(x))$  and similarly for  $\mathbf{qk}$ .
8. For sets  $\Gamma$  of formulae  $\mathbf{k}(\Gamma) := \bigcup \{\mathbf{k}(A) : A \in \Gamma\}$ .

For example  $\mathbf{qk}(\exists x \in a A(x)) = \{a\} \cup \mathbf{qk}(A(x))$  if  $a \in L_{\rho_0}$ .

**Definition 5.5** For  $a \in L_{\rho_0} \cup \{L_{\rho_0}\}$ ,  $\mathbf{rk}_L(a)$  denotes the  $L$ -rank of  $a$ .

$$\mathbf{rk}_L(a) := \begin{cases} \min\{\alpha \in \text{Ord} : a \in L_{\alpha+1}\} & a \in L_{\rho_0} \\ \rho_0 & a = L_{\rho_0} \end{cases}$$

**Definition 5.6** 1.  $A \in \Delta_0$  iff there exists a  $\Delta_0$ -formula  $\theta[\vec{x}]$  in the language  $\{\in\}$  and terms  $\vec{t}$  in  $\mathcal{L}_c$  such that  $A \equiv \theta[\vec{t}]$ . This means that  $A$  is bounded, and the predicates  $P, P_{\rho_0}$  do not occur in  $A$ .

2. Putting  $\Sigma_0 := \Pi_0 := \Delta_0$ , the classes  $\Sigma_m$  and  $\Pi_m$  of formulae in the language  $\mathcal{L}_c$  are defined as usual, where by definition  $\Sigma_m \cup \Pi_m \subset \Sigma_{m+1} \cap \Pi_{m+1}$ .

Each formula in  $\Sigma_m \cup \Pi_m$  is in prenex normal form with alternating unbounded quantifiers and  $\Delta_0$ -matrix.



3. The set  $\Sigma^{\Sigma_n}(\lambda)$  of sentences is defined recursively as follows. Let  $\{a, b, c\} \subset L_{\rho_0}$  and  $d \in L_{\rho_0} \cup \{L_{\rho_0}\}$ .
  - (a) Each  $\Sigma_n$ -sentence is in  $\Sigma^{\Sigma_n}(\lambda)$ .
  - (b) Each literal including  $Reg(a), P(a, b, c), P_{I,n}(a)$  and its negation is in  $\Sigma^{\Sigma_n}(\lambda)$ .
  - (c)  $\Sigma^{\Sigma_n}(\lambda)$  is closed under propositional connectives  $\vee, \wedge$ .
  - (d) Suppose  $\forall x \in d A(x) \notin \Delta_0$ . Then  $\forall x \in d A(x) \in \Sigma^{\Sigma_n}(\lambda)$  iff  $A(\emptyset) \in \Sigma^{\Sigma_n}(\lambda)$  and  $\text{rk}_L(d) < \lambda$ .
  - (e) Suppose  $\exists x \in d A(x) \notin \Delta_0$ . Then  $\exists x \in d A(x) \in \Sigma^{\Sigma_n}(\lambda)$  iff  $A(\emptyset) \in \Sigma^{\Sigma_n}(\lambda)$  and  $\text{rk}_L(d) \leq \lambda$ .
4. For a  $\Sigma_1$ -formula  $A(x)$ ,  $\exists x \in P_{\rho_0} A(x)$  is a  $\Sigma_1(P_{\rho_0})$ -formula.

Note that the predicates  $P, P_{\rho_0}$  do not occur in  $\Sigma_m$ -formulae.

**Definition 5.7** Let us extend the domain  $\text{dom}(F_x) = \text{Hull}(x)$  of the Mostowski collapse to formulae.

$$\text{dom}(F_x) = \{A \in \Sigma_1 \cup \Pi_1 : k(A) \subset \text{Hull}(x)\}.$$

For  $A \in \text{dom}(F_x)$ ,  $F_x''A$  denotes the result of replacing each constant  $c \in L_{\rho_0}$  by  $F_x(c)$ , each unbounded existential quantifier  $\exists z \in L_{\rho_0}$  by  $\exists z \in L_{F_x(\rho_0)}$ , and each unbounded universal quantifier  $\forall z \in L_{\rho_0}$  by  $\forall z \in L_{F_x(\rho_0)}$ .

For sequent, i.e., finite set of sentences  $\Gamma \subset \text{dom}(F_x)$ , put  $F_x''\Gamma = \{F_x''A : A \in \Gamma\}$ .

The assignment of disjunctions  $A \simeq \bigvee (A_\iota)_{\iota \in J}$  or conjunctions  $A \simeq \bigwedge (A_\iota)_{\iota \in J}$  to sentences  $A$  is defined as in [3] *except* for  $\Sigma_1 \cup \Pi_1$ -sentences.

**Definition 5.8** 1. If  $M$  is one of the literals  $a \in b, a \notin b$ , then for  $J := 0$

$$M \simeq \begin{cases} \bigvee (A_\iota)_{\iota \in J} & \text{if } M \text{ is false (in } L) \\ \bigwedge (A_\iota)_{\iota \in J} & \text{if } M \text{ is true} \end{cases}$$

2.  $(A_0 \vee A_1) \simeq \bigvee (A_\iota)_{\iota \in J}$  and  $(A_0 \wedge A_1) \simeq \bigwedge (A_\iota)_{\iota \in J}$  for  $J := 2$ .

3.  $P(b, c) \simeq \bigvee (0 \notin 0)_{\iota \in J}$  and  $\neg P(b, c) \simeq \bigwedge (0 \in 0)_{\iota \in J}$  with

$$J := \begin{cases} 1 & \text{if } \exists \alpha [b = \Psi_{\omega_1}(\alpha) \& c = F_{b \cup \{\omega_1\}}(\rho_0)] \\ 0 & \text{otherwise} \end{cases}.$$

4.  $P_{\rho_0}(a) \simeq \bigvee (0 \notin 0)_{\iota \in J}$  and  $\neg P_{\rho_0}(a) \simeq \bigwedge (0 \in 0)_{\iota \in J}$  with

$$J := \begin{cases} 1 & \text{if } \exists \alpha [a = \Psi_{\rho_0}(\alpha)] \\ 0 & \text{otherwise} \end{cases}.$$

5. Let  $\exists z \in b\theta[z] \in \Sigma_0$  for  $b \in L_{\rho_0} \cup \{L_{\rho_0}\}$ . Then for the set

$$d := \mu z \in b\theta[z] := \begin{cases} \min_{<_L} \{d : d \in b \wedge \theta[d]\} & \text{if } \exists z \in b\theta[z] \\ 0 & \text{otherwise} \end{cases}$$

with a canonical well ordering  $<_L$  on  $L$ , and  $J = \{d\}$

$$\begin{aligned} \exists z \in b\theta[z] & \simeq \bigvee (d \in b \wedge \theta[d])_{d \in J} \\ \forall z \in b \neg \theta[z] & \simeq \bigwedge (d \in b \rightarrow \neg \theta[d])_{d \in J} \end{aligned}$$

where  $d \in b$  denotes a true literal, e.g.,  $d \notin d$  when  $b = L_{\rho_0}$ .

6. For a  $\Sigma_1(P_{\rho_0})$ -sentence  $\exists x \in P_{\rho_0} A(x)$ ,

$$\begin{aligned} \exists x \in P_{\rho_0} A(x) & \simeq \bigvee (A(a))_{a \in J} \\ \forall x \in P_{\rho_0} \neg A(x) & \simeq \bigwedge (\neg A(a))_{a \in J} \\ \text{with } J & = \{a : \exists \alpha (a = \Psi_{\rho_0}(\alpha))\} \end{aligned}$$

7. Otherwise set for  $a \in L_{\rho_0} \cup \{L_{\rho_0}\}$  and  $J := \{b : b \in a\}$

$$\exists x \in a A(x) \simeq \bigvee (A(b))_{b \in J} \text{ and } \forall x \in a A(x) \simeq \bigwedge (A(b))_{b \in J}.$$

The rank  $\text{rk}(A)$  of sentences  $A$  is defined by recursion on the number of symbols occurring in  $A$ .

**Definition 5.9** 1.  $\text{rk}(\neg A) := \text{rk}(A)$ .

2.  $\text{rk}(a \in b) := 0$ .

3.  $\text{rk}(P(b, c)) := \text{rk}(P_{\rho_0}(a)) := 1$ .

4.  $\text{rk}(A_0 \vee A_1) := \max\{\text{rk}(A_0), \text{rk}(A_1)\} + 1$ .

5.  $\text{rk}(\exists x \in a A(x)) := \max\{\omega\alpha, \text{rk}(A(\emptyset)) + 1\}$  for  $\alpha = \text{rk}_L(a)$ .

6.  $\text{rk}(\exists x \in P_{\rho_0} A(x)) = \rho_0$ .

**Proposition 5.10** Let  $A \simeq \bigvee (A_\iota)_{\iota \in J}$  or  $A \simeq \bigwedge (A_\iota)_{\iota \in J}$ .

1.  $\forall \iota \in J (\mathbf{k}(A_\iota) \subset \mathbf{k}(A) \cup \{\iota\})$ .

2.  $A \in \Sigma^{\Sigma_n}(\lambda) \Rightarrow \forall \iota \in J (A_\iota \in \Sigma^{\Sigma_n}(\lambda))$ .

3. For an ordinal  $\lambda \leq \rho_0$  with  $\omega\lambda = \lambda$ ,  $\text{rk}(A) < \lambda \Rightarrow A \in \Sigma^{\Sigma_n}(\lambda)$ .

4.  $\text{rk}(A) < \rho_0 + \omega$ .

5.  $\text{rk}(A) \in \{\omega \text{rk}_L(a) + i : a \in \mathbf{qk}(A) \cup \{\rho_0\}, i \in \omega\} \subset \text{Hull}(\mathbf{k}(A))$ .

6.  $\forall \iota \in J (\text{rk}(A_\iota) < \text{rk}(A))$ .

### 5.3 Operator controlled derivations

In the remaining parts of this section we work in the intuitionistic fixed point theory  $\text{FiX}^i(T_1)$ .

Sequents are finite sets of sentences, and inference rules are formulated in one-sided sequent calculus. In what follows by an operator we mean an  $\mathcal{H}_\gamma[\Theta]$  for a finite set  $\Theta$  of sets.

**Definition 5.11** Define a relation  $(\mathcal{H}, \kappa) \vdash_b^a \Gamma$  as follows.

$(\mathcal{H}, \kappa) \vdash_b^a \Gamma$  holds if

$$\{a\} \cup \mathbf{k}(\Gamma) \subset \mathcal{H} := \mathcal{H}(\emptyset) \quad (6)$$

and one of the following cases holds:

( $\vee$ )  $A \simeq \bigvee \{A_\iota : \iota \in J\}$ ,  $A \in \Gamma$  and there exist  $\iota \in J$  and  $a(\iota) < a$  such that

$$\text{rk}_L(\iota) < \kappa \Rightarrow \text{rk}_L(\iota) < a \quad (7)$$

and  $(\mathcal{H}, \kappa) \vdash_b^{a(\iota)} \Gamma, A_\iota$ .

( $\wedge$ )  $A \simeq \bigwedge \{A_\iota : \iota \in J\}$ ,  $A \in \Gamma$  and for every  $\iota \in J$  there exists an  $a(\iota) < a$  such that  $(\mathcal{H}[\{\iota\}], \kappa) \vdash_b^{a(\iota)} \Gamma, A_\iota$ .

(*cut*) There exist  $a_0 < a$  and  $C$  such that  $\text{rk}(C) < b$  and  $(\mathcal{H}, \kappa) \vdash_b^{a_0} \Gamma, \neg C$  and  $(\mathcal{H}, \kappa) \vdash_b^{a_0} C, \Gamma$ .

(**P**) There exists  $\alpha < \omega_1$  such that  $(\exists x, y < \omega_1 [\alpha < x \wedge P(x, y)]) \in \Gamma$ .

(**F** $_{x \cup \{\omega_1\}}$ )  $x = \Psi_{\omega_1}(\beta) \in \mathcal{H}$  for a  $\beta$  and there exist  $a_0 < a$ ,  $\Gamma_0 \subset \Sigma_1$  and  $\Lambda$  such that  $\mathbf{k}(\Gamma_0) \subset \text{Hull}(x \cup \{\omega_1\})$ ,  $\Gamma = \Lambda \cup (F_{x \cup \{\omega_1\}} \text{"} \Gamma_0)$  and  $(\mathcal{H}, \kappa) \vdash_b^{a_0} \Lambda, \Gamma_0$ , where  $F_{x \cup \{\omega_1\}}$  denotes the Mostowski collapse  $F_{x \cup \{\omega_1\}} : \text{Hull}(x \cup \{\omega_1\}) \leftrightarrow L_{F_{x \cup \{\omega_1\}}(\rho_0)}$ .

(**P** $_{\rho_0}$ ) There exists  $\alpha < \rho_0$  such that  $(\exists x < \rho_0 [\alpha < x \wedge P_{\rho_0}(x)]) \in \Gamma$ .

(**F** $_x$ )  $x = \Psi_{\rho_0}(\beta) \in \mathcal{H}$  for a  $\beta$  and there exist  $a_0 < a$ ,  $\Gamma_0 \subset \Sigma_1$  and  $\Lambda$  such that  $\mathbf{k}(\Gamma_0) \subset \text{Hull}(x)$ ,  $\Gamma = \Lambda \cup (F_x \text{"} \Gamma_0)$  and  $(\mathcal{H}, \kappa) \vdash_b^{a_0} \Lambda, \Gamma_0$ , where  $F_x$  denotes the Mostowski collapse  $F_x : \text{Hull}(x) \leftrightarrow L_{F_x(\rho_0)}$ .

(*Ref*)  $b \geq \rho_0$ , and there exist an ordinal  $a_0 < a$ , a set  $c$  and a  $\Sigma_1(P_{\rho_0})$ -formula  $A(x)$  such that  $(\mathcal{H}, \kappa) \vdash_b^{a_0} \Gamma, \forall x \in c A(x)$  and  $(\mathcal{H}, \kappa) \vdash_b^{a_0} \forall y \exists x \in c \neg A^{(y)}(x), \Gamma$ , where for  $A(x) \equiv (\exists z \in P_{\rho_0} \exists w B(x))$  ( $B \in \Delta_0$ ),  $A^{(y)}(x) := (\exists z \in P_{\rho_0} \cap y \exists w \in y B)$ .

**Lemma 5.12** (Tautology) *If  $\mathbf{k}(\Gamma \cup \{A\}) \subset \mathcal{H}$  then  $(\mathcal{H}, \rho_0) \vdash_0^{2\text{rk}(A)} \Gamma, \neg A, A$ .*

**Lemma 5.13** *Let  $\text{rk}(\forall x \in b \varphi[x, c]) \leq \rho_0 + m$  for an  $m \geq 1$ , and  $\Theta_c = \{\neg \forall y (\forall x \in y \varphi[x, c] \rightarrow \varphi[y, c])\}$ . Then for any operator  $\mathcal{H}$ , and any  $a, c$ ,  $(\mathcal{H}[\{c, a\}], \rho_0) \vdash_{\rho_0+m+1}^{\rho_0+2m+2+2\text{rk}_L(a)} \Theta_c, \forall x \in a \varphi[x, c]$ .*

Let

$$\begin{aligned} (\mathcal{H}, \rho_0) \vdash_c^{\leq \alpha} \Gamma & :\Leftrightarrow \exists \beta < \alpha [(\mathcal{H}, \rho_0) \vdash_c^\beta \Gamma] \\ (\mathcal{H}, \rho_0) \vdash_c^{\leq \alpha} \Gamma & :\Leftrightarrow \exists d < c [(\mathcal{H}, \rho_0) \vdash_d^{\leq \alpha} \Gamma] \end{aligned}$$

**Lemma 5.14** *Let  $A$  be an axiom in  $T(\omega_1)$  except Foundation axiom schema and  $\Pi_1$ -Collection. Then  $(\mathcal{H}, \rho_0) \vdash_0^{\leq \rho_0 + \omega} A$  for any operator  $\mathcal{H} = \mathcal{H}_\gamma$ .*

**Lemma 5.15** (Embedding)

*If  $T(\omega_1) \vdash \Gamma[\vec{x}]$ , there are  $m, k < \omega$  such that for any  $\vec{a} \subset L_{\rho_0}$ ,  $(\mathcal{H}[\vec{a}], \rho_0) \vdash_{\rho_0+m}^{\rho_0 \cdot 2 + k} \Gamma[\vec{a}]$  for any operator  $\mathcal{H} = \mathcal{H}_\gamma$ .*

**Proof.**

By Lemma 5.13 we have  $(\mathcal{H}, \rho_0) \vdash_{\rho_0+m+1}^{\rho_0 \cdot 2} \forall u, z (\forall y (\forall x \in y \varphi[x, z] \rightarrow \varphi[y, z]) \rightarrow \varphi[u, z])$  for some  $m$ . By Lemmata 5.12 and 5.14 it remains to consider instances

$$\forall u \in a \exists v \forall w \theta \rightarrow \exists c \forall u \in a \exists v \in c \forall w \theta$$

of  $\Pi_1$ -Collection, where  $\theta \equiv \theta(u, v, w)$  is a  $\Delta_0$ -formula in the language  $\{\in\}$ .

First by Lemma 5.14 with axioms (3) and (4) we have

$$(\mathcal{H}, \rho_0) \vdash_{\rho_0+1}^{\rho_0+\omega} \forall w \theta(u, v, w) \leftrightarrow \exists x \in P_{\rho_0} \tau(x, u, v)$$

where  $\tau(x, u, v) \equiv [u, v \in L_x \wedge \forall w \in L_x \theta(u, v, w)]$ . Hence

$$(\mathcal{H}, \rho_0) \vdash_{\leq \rho_0 + \omega}^{\leq \rho_0 + \omega \cdot 2} \neg \forall u \in a \exists v \forall w \theta, \forall u \in a \exists x \in P_{\rho_0} \exists v \tau(x, u, v)$$

On the other hand we have by Lemma 5.12

$$(\mathcal{H}, \rho_0) \vdash_0^{\leq \rho_0 + \omega} \neg \exists c \forall u \in a \exists x \in P_{\rho_0} \cap c \exists v \in c \tau, \exists c \forall u \in a \exists x \in P_{\rho_0} \cap c \exists v \in c \tau$$

Hence by the inference (*Ref*) for the  $\Sigma_1(P_{\rho_0})$ -formula  $\exists x \in P_{\rho_0} \exists v \tau(x, u, v)$ , we obtain

$$(\mathcal{H}, \rho_0) \vdash_{\leq \rho_0 + \omega}^{\leq \rho_0 + \omega \cdot 2} \neg \forall u \in a \exists v \forall w \theta, \exists c \forall u \in a \exists x \in P_{\rho_0} \cap c \exists v \in c \tau$$

Therefore  $(\mathcal{H}, \rho_0) \vdash_{\leq \rho_0 + \omega}^{\rho_0 + \omega \cdot 2} \forall u \in a \exists v \forall w \theta \rightarrow \exists c \forall u \in a \exists v \in c \forall w \theta$ .  $\square$

In the following Lemma 5.16, note that  $\text{rk}(\exists x < \omega_1 \exists y < \omega_1 [\alpha < x \wedge P(x, y)]) = \omega_1 + 1$ , and  $\text{rk}(\exists x < \rho_0 [\alpha < x \wedge P_{\rho_0}(x)]) = \rho_0$ .

**Lemma 5.16** (Predicative Cut-elimination)

1. If  $(\mathcal{H}, \kappa) \vdash_{c+\omega^a}^b \Gamma \& [c, c + \omega^a [\cap \{\omega_1 + 1, \rho_0\} = \emptyset \& a \in \mathcal{H} \Rightarrow (\mathcal{H}, \kappa) \vdash_c^{\varphi ab} \Gamma]$ .
2. If  $(\mathcal{H}_\gamma, \kappa) \vdash_{\omega_1+2}^b \Gamma \& \gamma \in \mathcal{H}_\gamma \Rightarrow (\mathcal{H}_{\gamma+b}, \kappa) \vdash_{\omega_1+1}^{\omega^b} \Gamma$ .
3. If  $(\mathcal{H}_\gamma, \kappa) \vdash_{\rho_0+1}^b \Gamma \& \gamma \in \mathcal{H}_\gamma \Rightarrow (\mathcal{H}_{\gamma+b}, \kappa) \vdash_{\rho_0}^{\omega^b} \Gamma$ .

For a formula  $\exists x \in d A(x)$  and ordinals  $\lambda = \text{rk}_L(d), \alpha$ ,  $(\exists x \in d A)^{(\exists \lambda | \alpha)}$  denotes the result of restricting the outermost existential quantifier  $\exists x \in d$  to  $\exists x \in L_\alpha$ ,  $(\exists x \in d A)^{(\exists \lambda | \alpha)} \equiv (\exists x \in L_\alpha A)$ .

**Lemma 5.17** (Boundedness) *Let  $\lambda \in \{\omega_1, \rho_0\}$ ,  $C \equiv (\exists x \in d A) \in \Sigma^{\Sigma_2}(\lambda)$  and  $C \notin \{\exists x < \omega_1 \exists y < \omega_1 [\alpha < x \wedge P(x, y)] : \alpha < \omega_1\} \cup \{\exists x < \rho_0 [\alpha < x \wedge P_{\rho_0}(x)] : \alpha < \rho_0\}$ .*

$$1. (\mathcal{H}, \lambda) \vdash_c^a \Lambda, C \& a \leq b \in \mathcal{H} \cap \lambda \Rightarrow (\mathcal{H}, \lambda) \vdash_c^a \Lambda, C^{(\exists \lambda | b)}.$$

$$2. (\mathcal{H}, \kappa) \vdash_c^a \Lambda, \neg C \& b \in \mathcal{H} \cap \lambda \Rightarrow (\mathcal{H}, \kappa) \vdash_c^a \Lambda, \neg(C^{(\exists \lambda | b)}).$$

**Lemma 5.18** (Boundedness for  $\Sigma_1(P_{\rho_0})$ )

*Let  $C$  be a  $\Sigma_1(P_{\rho_0})$ -sentence. Then  $(\mathcal{H}, \rho_0) \vdash_c^a \Lambda, C \& a \leq b \in \mathcal{H} \cap \rho_0 \Rightarrow (\mathcal{H}, \rho_0) \vdash_c^a \Lambda, C^{(L_b)}$ .*

**Proof.**  $C^{L_b} \equiv (\exists z \in P_{\rho_0} \cap L_b \exists w \in L_b B)$  when  $C \equiv (\exists z \in P_{\rho_0} \exists w B)$  with a  $\Delta_0$ -formula  $B$ . The lemma is seen from (7).  $\square$

## 5.4 Collapsing derivations

In this subsection derivations of  $\Sigma^{\Sigma_2}(\omega_1)$  sentences are shown to be collapsed to derivations with heights and cut ranks  $< \omega_1$ .

**Lemma 5.19** (Collapsing below  $\omega_1$ )

*Suppose  $\gamma \in \mathcal{H}_\gamma[\Theta]$  with  $\Theta \subset \mathcal{H}_\gamma(\Psi_{\omega_1}(\gamma))$ , and  $\Gamma \subset \Sigma^{\Sigma_2}(\omega_1)$ . Then for  $b = \Psi_{\omega_1}(\gamma + \omega^{\omega_1+a})$ ,*

$$(\mathcal{H}_\gamma[\Theta], \omega_1) \vdash_{\omega_1+1}^a \Gamma \Rightarrow (\mathcal{H}_{\gamma+\omega^{\omega_1+a}+1}[\Theta], \omega_1) \vdash_b^b \Gamma.$$

**Lemma 5.20** (Collapsing below  $\rho_0$ )

*Suppose  $\gamma \in \mathcal{H}_\gamma[\Theta]$  with  $\Theta \subset \mathcal{H}_\gamma(\Psi_{\rho_0}(\gamma))$ , and  $\Gamma \subset \Sigma^{\Sigma_2}(\rho_0) \cup \Sigma_1(P_{\rho_0})$ . Then for  $\hat{a} = \gamma + \omega^{\rho_0+a}$*

$$(\mathcal{H}_\gamma[\Theta], \rho_0) \vdash_{\rho_0}^a \Gamma \Rightarrow (\mathcal{H}_{\hat{a}+1}[\Theta], \rho_0) \vdash_{\Psi_{\rho_0}(\hat{a})}^{\Psi_{\rho_0}(\hat{a})} \Gamma.$$

**Proof** by induction on  $a$ , cf. Lemma 5.1.

First note that  $\Psi_{\rho_0}(\hat{a}) \in \mathcal{H}_{\hat{a}+1}[\Theta]$  since  $\hat{a} = \gamma + \omega^{\rho_0+a} \in \mathcal{H}_\gamma[\Theta] \subset \mathcal{H}_{\hat{a}+1}[\Theta]$  by the assumption,  $\{\gamma, a\} \subset \mathcal{H}_\gamma[\Theta]$ .

Assume  $(\mathcal{H}_\gamma[\Theta][\Theta_0], \rho_0) \vdash_{\rho_0}^{a_0} \Gamma_0$  with  $\Theta_0 \subset \mathcal{H}_\gamma(\Psi_{\rho_0}(\gamma))$ . Then by  $\gamma \leq \hat{a}$ , we have  $\hat{a}_0 \in \mathcal{H}_\gamma[\Theta][\Theta_0] \subset \mathcal{H}_\gamma(\Psi_{\rho_0}(\gamma)) \subset \mathcal{H}_{\hat{a}}(\Psi_{\rho_0}(\hat{a}))$ . This yields that

$$a_0 < a \Rightarrow \Psi_{\rho_0}(\hat{a}_0) < \Psi_{\rho_0}(\hat{a})$$

Second observe that  $k(\Gamma) \subset \mathcal{H}_\gamma[\Theta] \subset \mathcal{H}_{\hat{a}+1}[\Theta]$  by  $\gamma \leq \hat{a} + 1$ .

Third we have

$$k(\Gamma) \subset \mathcal{H}_\gamma(\Psi_{\rho_0}(\gamma))$$

When  $\Gamma$  is one of axioms  $(\mathbf{P})$  and  $(\mathbf{P}_{\rho_0})$ , there is nothing to show.

Consider the case when the last inference is a  $(Ref)$ .

$$\frac{(\mathcal{H}_\gamma[\Theta], \rho_0) \vdash_{\rho_0}^{a_0} \Gamma, \forall x \in c A(x) \quad (\mathcal{H}_\gamma[\Theta], \rho_0) \vdash_{\rho_0}^{a_0} \forall y \exists x \in c \neg A^{(y)}(x), \Gamma}{(\mathcal{H}_\gamma[\Theta], \rho_0) \vdash_{\rho_0}^a \Gamma} (Ref)$$

where  $a_0 < a$  and  $A(x) \equiv (\exists z \in P_{\rho_0} \exists w B(x))$  is a  $\Sigma_1(P_{\rho_0})$ -formula with a  $\Delta_0$ -formula  $B$ .

For each  $d \in c$  we have by Inversion

$$(\mathcal{H}_\gamma[\Theta \cup \{d\}], \rho_0) \vdash_{\rho_0}^{a_0} \Gamma, A(d)$$

where  $c \in \mathcal{H}_\gamma(\Psi_{\rho_0}(\gamma))$ . Hence  $\text{rk}_L(d) < \text{rk}_L(c) \in \mathcal{H}_\gamma(\Psi_{\rho_0}(\gamma)) \cap \rho_0 \subset \Psi_{\rho_0}(\gamma)$ , and  $\text{rk}_L(d) < \Psi_{\rho_0}(\gamma)$ . Therefore  $d \in \mathcal{H}_\gamma(\Psi_{\rho_0}(\gamma))$ . By IH we have for  $\hat{a}_0 = \gamma + \omega^{\rho_0 + a_0}$  and  $\beta_0 = \Psi_{\rho_0}(\hat{a}_0) \in \mathcal{H}_{\hat{a}_0+1}[\Theta]$

$$(\mathcal{H}_{\hat{a}_0+1}[\Theta \cup \{d\}], \rho_0) \vdash_{\beta_0}^{\beta_0} \Gamma, A(d)$$

Boundedness lemma 5.18 yields

$$(\mathcal{H}_{\hat{a}_0+1}[\Theta \cup \{d\}], \rho_0) \vdash_{\beta_0}^{\beta_0} \Gamma, A^{(L_{\beta_0})}(d)$$

Since  $d \in c$  is arbitrary, we obtain by  $(\bigwedge)$

$$(\mathcal{H}_{\hat{a}_0+1}[\Theta], \rho_0) \vdash_{\beta_0}^{\beta_0+1} \Gamma, \forall x \in c A^{(L_{\beta_0})}(x) \quad (8)$$

On the other hand we have by Inversion for  $L_{\beta_0} \in \mathcal{H}_{\hat{a}_0+1}[\Theta]$

$$(\mathcal{H}_{\hat{a}_0+1}[\Theta], \rho_0) \vdash_{\rho_0}^{a_0} \exists x \in c \neg A^{(L_{\beta_0})}(x), \Gamma$$

Since  $\exists x \in c \neg A^{(L_{\beta_0})}(x) \in \Sigma^{\Sigma_2}(\rho_0)$ , IH yields for  $\hat{a}_1 = \hat{a}_0 + 1 + \omega^{\rho_0 + a_0} = \gamma + \omega^{\rho_0 + a_0} \cdot 2$  and  $\beta_1 = \Psi_{\rho_0}(\hat{a}_1)$

$$(\mathcal{H}_{\hat{a}_1+1}[\Theta], \rho_0) \vdash_{\beta_1}^{\beta_1} \exists x \in c \neg A^{(L_{\beta_0})}(x), \Gamma \quad (9)$$

We have  $\text{rk}(\forall x \in c A^{(L_{\beta_0})}(x)) \in \text{Hull}(\text{k}(\forall x \in c A^{(L_{\beta_0})}(x))) \cap \rho_0 \subset \mathcal{H}_{\hat{a}_0+1}[\Theta] \cap \rho_0 \subset \mathcal{H}_{\hat{a}_0+1}(\Psi_{\rho_0}(\gamma)) \cap \rho_0 \subset \Psi_{\rho_0}(\hat{a})$  by Proposition 5.10.5.

By a *cut* with (8) and (9) we obtain with  $\Psi_{\rho_0}(\hat{a}) > \beta_1 > \beta_0$

$$(\mathcal{H}_{\hat{a}+1}[\Theta], \rho_0) \vdash_{\Psi_{\rho_0}(\hat{a})}^{\Psi_{\rho_0}(\hat{a})} \Gamma$$

Other case as seen as in [1].  $\square$

## 6 Proof of Theorem 1.1

For a sentence  $\exists x \in L_{\omega_1} \varphi$  with a  $\Sigma_2$ -formula  $\varphi$  in the language  $\{\in, \omega_1\}$ , assume  $T_1 \vdash \exists x \in L_{\omega_1} \varphi$ . Then by Lemmata 3.2 and 5.15, pick an  $m > 0$  such that the

fact  $(\mathcal{H}_0, \rho_0) \vdash_{\rho_0+m}^{\rho_0 \cdot 2+m} \exists x \in L_{\omega_1} \varphi$  is provable in  $\text{FiX}^i(T_1)$ . In what follows work in  $\text{FiX}^i(T_1)$ . Predicative Cut Elimination 5.16.1 and 5.16.3 yields

$$(\mathcal{H}_\gamma, \rho_0) \vdash_{\rho_0}^a \exists x \in L_{\omega_1} \varphi$$

for  $\gamma = \omega_{m-1}(\rho_0 \cdot 2 + m)$  and  $a = \omega_m(\rho_0 \cdot 2 + m)$ . Then Collapsing below  $\rho_0$  5.20 yields

$$(\mathcal{H}_{\omega_{m+1}(\rho_0 \cdot 2+m)+1}, \rho_0) \vdash_{\beta}^{\beta} \exists x \in L_{\omega_1} \varphi$$

for  $\gamma + \omega^{\rho_0+a} = \omega_{m+1}(\rho_0 \cdot 2 + m)$  and  $\beta = \Psi_{\rho_0}(\omega_{m+1}(\rho_0 \cdot 2 + m))$ . Predicative Cut Elimination 5.16.1 and 5.16.2 yields

$$(\mathcal{H}_{\omega_{m+1}(\rho_2 \cdot 2+m)+\varphi\beta\beta}, \rho_0) \vdash_{\omega_1+1}^{\varphi\beta\beta} \exists x \in L_{\omega_1} \varphi$$

for  $\omega_1 + 2 + \omega^{\beta} = \beta$  and  $\omega^{\varphi\beta\beta} = \varphi\beta\beta$ . A fortiori,

$$(\mathcal{H}_{\omega_{m+1}(\rho_2 \cdot 2+m)+\varphi\beta\beta}, \omega_1) \vdash_{\omega_1+1}^{\varphi\beta\beta} \exists x \in L_{\omega_1} \varphi$$

Then Collapsing below  $\omega_1$  5.19 yields

$$(\mathcal{H}_{\omega_{m+1}(\rho_2 \cdot 2+m)+(\varphi\beta\beta \cdot 2)+1}, \omega_1) \vdash_{\delta}^{\delta} \exists x \in L_{\omega_1} \varphi$$

for  $\omega_{m+1}(\rho_2 \cdot 2 + m) + \varphi\beta\beta + \omega^{\omega_1+\varphi\beta\beta} + 1 = \omega_{m+1}(\rho_2 \cdot 2 + m) + (\varphi\beta\beta \cdot 2) + 1$  and  $\delta = \Psi_{\omega_1}(\omega_{m+1}(\rho_2 \cdot 2 + m) + (\varphi\beta\beta \cdot 2))$ .

Boundedness 5.17.1 yields for  $\delta < \Psi_{\omega_1}(\omega_n(\rho_0 + 1))$  with  $n = m + 2$

$$(\mathcal{H}_{\omega_n(\rho_0+1)+1}, \omega_1) \vdash_{\delta}^{\delta} \exists x \in L_{\Psi_{\omega_1}(\omega_n(\rho_0+1))} \varphi$$

We see then by transfinite induction up to the countable ordinal  $\delta$  that inference rules in the controlled derivation of  $\exists x \in L_{\Psi_{\omega_1}(\omega_n(\rho_0+1))} \varphi$  with cut rank  $< \omega_1$  are  $(\vee)$ ,  $(\wedge)$ ,  $(cut)$ , and  $(\mathbf{F}_{x \cup \{\omega_1\}})$ , and since these inference rules are truth-preserving, we conclude again by transfinite induction up to  $\delta$  that  $\exists x \in L_{\Psi_{\omega_1}(\omega_n(\rho_0+1))} \varphi$  is true.

Since the whole proof is formalizable in  $\text{FiX}^i(T_1)$ , we conclude  $\text{FiX}^i(T_1) \vdash \exists x \in L_{\Psi_{\omega_1}(\omega_n(\rho_0+1))} \varphi$ . Finally Theorem 5.2 yields  $T_1 \vdash \exists x \in L_{\Psi_{\omega_1}(\omega_n(\rho_0+1))} \varphi$ . This completes a proof of Theorem 1.1.

## References

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- [3] W. Buchholz, A simplified version of local predicativity, P. H. G. Aczel, H. Simmons and S. S. Wainer(eds.), *Proof Theory*, Cambridge UP, 1992, pp. 115-147.